Dynamic Programming for Boolean decisions

Lecture 07.03 by Marina Barsky

Mix tapes

Subset sum

Mixtapes

- The mixtape problem is inspired by making <u>musical mixes</u> on cassettes (and later CDs)
- Given a set of songs with their durations, the question is whether these songs can be divided into 2 subsets where the total duration of each subset is the same



Réducteur de Souffle EREKY THE

Mixtape problem

Input: The durations of *n* songs $d_1, d_2, ..., d_n$ in minutes (integers).

Output: Yes, if the songs can be divided into two groups, such that each group has the same total duration. *No*, otherwise.

Sample problem instance

	1	2	3	4	5	6	7
min	3	2	3	2	2	5	3

A total of 7 songs given with their durations.

The output for this instance of the problem is 'yes'. This is because the songs can be divided into two groups that both have a total duration of 10 min.

3	2	3	2	2	5	3
1	2	3	4	5	6	7

The blue and red songs have the same total duration of 10 min.

DP solution: brainstorming

What would help us to know if a set of numbers can be divided into 2 subsets with equal sums?

How can we find out if there is a subset with a given sum?

Uwhat are optimal subproblems?

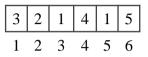
Think of an optimal solution to a subset sum

☐ If there is a subset with total duration D, and it contains song i, then there also should be a subset with duration D-d_i



As always, we can start by checking if all possible durations from 1 to D can be obtained from a current set, and we will reuse this knowledge to obtain an answer for duration D

Example

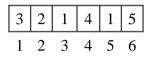


□ First, we compute total duration: 3+2+1+4+1+5 = 16

- The task becomes to find out if there is a subset that sums up to 16/2 = 8
- We will try methodically to fit each song into the solution, checking if the following total durations are possible: 0,1,2,3,4,5,6,7 and finally 8.

The check will produce a boolean value: Y(True) or N(False)

Create DP table



Л

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)									
d ₂ (2)									
d ₃ (1)									
d ₄ (4)									
d ₅ (1)									
d ₆ (5)									

Base condition

Is it possible to create a subset with a total duration **0**? Yes, just do not take any song.

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	т								
d ₂ (2)	т								
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	т	F							
d ₂ (2)	т								
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	т	F	F						
d ₂ (2)	т								
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	т	F	F	т					
d ₂ (2)	т								
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	т	F	F	Т	F	F	F	F	F
d ₂ (2)	т								
d ₃ (1)	Т								
d ₄ (4)	т								
d ₅ (1)	т								
d ₆ (5)	Т								

Using only song 1 (duration 3) and/or song 2(duration 2), it is still not possible to create a subset with total duration 1.

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F							
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	т								
d ₆ (5)	Т								

Using only song 1 (duration 3) and song 2(duration 2), it is not possible to create a subset with total duration 1, but it is possible to create a subset with total duration 2

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	Т	F	Т						
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

Using only d1 and d2, can we have a subset with total duration 3?

Yes, we already know that we can do it even without d2

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т					
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

How do we check if a subset sum 4 is possible? We know that it was False when we used d1 only, so if we use d2, then we need to check if a subset of (4-2) was possible. It was not.

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F				
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

To check for d=5, take item d2(2) and see if duration 5-2 was possible with the previous item(s)

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	т	F	F	F
d ₃ (1)	Т								
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

Considering songs 1,2, and 3

1,2,3 are possible.

What about 4? Current item d3 has duration 1. Is it possible to have a duration (4-1) with the other 2 items? Yes

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	т				
d ₄ (4)	Т								
d ₅ (1)	Т								
d ₆ (5)	Т								

Considering songs 1,2, and 3

Using only items d1, d2, d3 we get the following boolean answers.

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	т	F	F	F
d ₃ (1)	т	Т	Т	Т	Т	Т	Т	F	F
d ₄ (4)	т								
d ₅ (1)	т								
d ₆ (5)	т								

Considering songs 1,2,3,4

For d4(4) we do not even need to consider this item for durations 1,2,3,4,5,6 - we could make these subsets even without item d4(4).

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	Т	Т	Т	F	F
d ₄ (4)	Т	Т	Т	Т	Т	Т	Т		
d ₅ (1)	Т								
d ₆ (5)	Т								

Considering songs 1,2,3,4

What about 7? Fit d4(4) and see if (7-4) was True.

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	Т	Т	Т	F	F
d ₄ (4)	т	Т	Т	Т	Т	Т	Т	Т	
d ₅ (1)	Т								
d ₆ (5)	Т								

Considering songs 1,2,3,4

Same holds for total duration 8.

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	Т	Т	Т	F	F
d ₄ (4)	Т	Т	Т	Т	Т	Т	Т	Т	Т
d ₅ (1)	Т								
d ₆ (5)	Т								

Is total duration 8 possible?

At this point we can stop. We know that it is possible to form a subset with a total duration 8 even using only the first 4 items. But what is this subset?

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	т	Т	Т	Т	Т	F	F
d ₄ (4)	Т	Т	Т	Т	Т	Т	Т	Т	Т
d ₅ (1)	т								
d ₆ (5)	Т								

Recovering the subset with sum 8: trace back

The subset clearly includes item d4 - without it 8 was not possible

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	Т	Т	Т	F	F
d ₄ (4)	Т	Т	Т	Т	Т	Т	Т	Т	Т
d ₅ (1)	т								
d ₆ (5)	Т								

Recovering the subset with sum 8: trace back

If it includes item 4, we need to look at total duration (8-4). This one only became True when we added item d3.

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	т	Т	Т	F	F
d ₄ (4)	Т	Т	Т	Т	Т	Т	Т	Т	Т
d ₅ (1)	Т								
d ₆ (5)	Т								

Recovering the subset with sum 8: trace back

If the solution includes item d3(1), we need to look at total duration (4-1). This one is True because a previous item produced True. This item was d1(3)

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	т	Т	Т	F	F
d ₄ (4)	Т	Т	Т	Т	Т	Т	Т	Т	Т
d ₅ (1)	Т								
d ₆ (5)	Т								

Answer: Yes, it is possible to create 2 subsets with equal total duration

d1(3) + d3(1) + d4(4) = d2(2) + d5(1) + d6(5)

3 + 1 + 4 = 2 + 1 + 5

Total duration→	0	1	2	3	4	5	6	7	8
d ₁ (3)	Т	F	F	Т	F	F	F	F	F
d ₂ (2)	Т	F	Т	Т	F	Т	F	F	F
d ₃ (1)	Т	Т	Т	Т	т	Т	Т	F	F
d ₄ (4)	Т	Т	Т	Т	Т	Т	Т	Т	Т
d ₅ (1)	Т								
d ₆ (5)	Т								

Game of Rocks

Optimal game strategy

Game: 1-2 rocks

- 2 players
- 2 piles of rocks:



- with *n* and *m* rocks respectively
- Each turn, one player may take either 1 rock (from either pile) or 2 rocks (one from each pile)
- Once the rocks are taken, they are removed from play
- The player that takes the last rock wins

Winning strategy with DP

- □ To find the winning strategy for the m + n game, we first construct an mxn table *R*.
- □ If Player 1 can always win the n + m game, then we would say R(n, m) = W, but if Player 1 has no winning strategy against a player that always makes the right moves, we would write R(n, m) = L.
- \Box Computing R(n, m) for arbitrary n and m seems difficult, but we can build on smaller values.

DP table for game outcomes

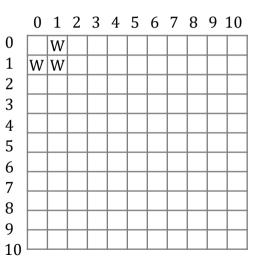
Construct an mxn table *R*.

□ Example: let m=n=10.

0 1 2 3 4 5 6 7 8 9 10

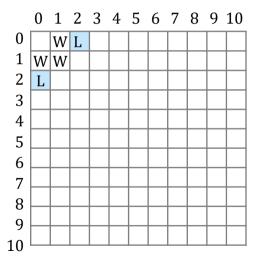
Simple subproblems first

- Notably R(0, 1), R(1, 0), and R(1, 1), are clearly winning propositions for Player 1 since with a single move Player 1 can win.
- Thus, we fill in entries (1, 1), (0, 1), and (1, 0) as W.



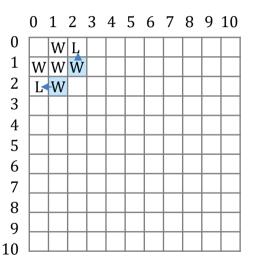
Solve larger subproblems based on solutions to the smaller problems

- In the (2, 0) case, the only move that Player 1 can make leads to the (1, 0) case that, as we already know, is a winning position for his opponent.
- A similar analysis applies to the (0, 2) case.



Solve larger subproblems based on solutions to the smaller problems

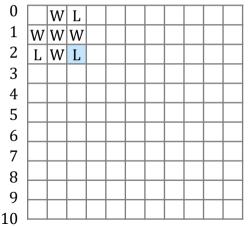
- In the (2, 1) case, Player 1 can make 3 different moves that lead respectively to the games of (1, 1), (2, 0), or (1, 0).
- One of these cases, (2, 0), leads to a losing position for his opponent and therefore (2, 1) is a winning position.
- The case (1, 2) is symmetric to (2, 1)



Solve larger subproblems based on solutions to the smaller problems

- In the (2, 2) case,
 Player 1 can make
 three different moves
 that lead to entries (2,
 1), (1, 2), and (1, 1).
- All of these entries are winning positions for his opponent and therefore R(2, 2) = L.

0 1 2 3 4 5 6 7 8 9 10



Fill DP table with game outcomes

- We can proceed filling in R in this way by noticing that for the entry (i, j) to be L, all the entries above, diagonally to the left, and directly to the left, must be W.
- These entries:

((i -1, j), (i -1, j -1), (i, j -1))correspond to the three possible moves that Player 1 can make.

0		W	I.	W	I.	W	I.	W	L	W	L
1	w				_	_	_	_	_	_	W
2	L		L					_	_	W	
3	w	_	_	_	_	_	_	_	_		W
4	L		L			_	_	_		W	
5	W	W	W	W	W	W	W	W	W	W	W
6	L	W	L	W	L	W	L	W	L	W	L
7	W	W	W	W	W	W	W	W	W	W	W
8	L	W	L	W	L	W	L	W	L	W	L
9	W	W	W	W	W	W	W	W	W	W	W
0	L	W	L	W	L	W	L	W	L	W	L

Rocks: winning strategy

□ The *Rocks* algorithm determines if Player 1 wins or loses.

□ If Player 1 wins in an *n+m* game, *Rocks* returns W. If Player 1 loses, *Rocks* returns L.

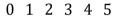
We introduce an artificial initial condition, R(0, 0)
 = L to simplify the pseudocode.

Algorithm *Rocks*(*n*, *m*)

```
R[0, 0] \leftarrow L
for i from 1 to n:
                                     # initialize rows
  if R[i - 1, 0] = W:
    R[i, 0] \leftarrow L
  else:
    R[i, 0] \leftarrow W
for j from 1 to m:
                                     # initialize columns
  if R[0, i - 1] = W:
    R[0, i] \leftarrow L
  else:
     R[0, i] \leftarrow W
for i from 1 to n:
  for j from 1 to m:
                                     # fill DP table
      if R[i - 1, j - 1] = W and R[i, j - 1] = W and R[i - 1, j] = W:
        R[i, i] \leftarrow L
      else:
         R[i, i] \leftarrow W
return R[n, m]
```

Using DP table for best strategy or game AI

- We can use the DP table to always play the winning strategy.
- If R(n,m) = W, and Player 1 starts first, he can always win: by taking the number of rocks which lead to the losing position of our opponent.
- If R(n,m) = L, then Player 1 can only hope that Player 2 does not use the same table, and makes a mistake.



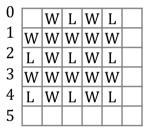
0		W	L	W	L	W
1	W	W	W	W	W	W
2	L	W	L	W	L	W
3	W	W	W	W	W	W
4	L	W	L	W	L	W
5	W	W	W	W	W	W

0 1 2 3 4 5

0		W	L	W	L	W
1	W	W	W	W	W	W
2	L	W	L	W	L	W
3	W	W	W	W	W	W
4	L	W	L	W	L	W
5	W	W	W	W	W	W

Player 1 takes (1,1).

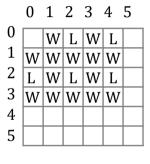




Player 1 takes (1,1).

No matter what Player 2 does, it leads to the winning state of Player 1.

Say, Player 2 takes (1,0)

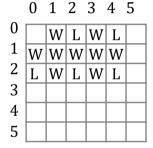


Player 2 takes (1,0)

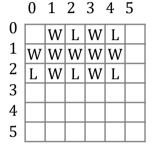
 $0 \ 1 \ 2 \ 3 \ 4 \ 5$

0		W	L	W	L	
	W	W	W	W	W	
2	L	W	L	W	L	
3	W	W	W	W	W	
4						
5						

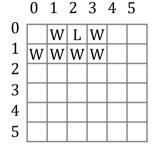
Player 1 should take (1,0).



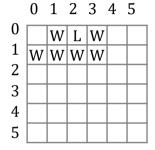
Player 1 takes (1,0).



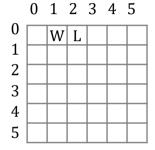
Player 2 takes (1,1).



Player 2 takes (1,1).



Player 1 should take (1,1).



Player 1 takes (1,1).

At this point the victory for Player 1 is guaranteed.

Identifying patterns

- A faster algorithm relies on the simple pattern in R, and checks if n and m are both even, in which case the player 1 loses.
- However, though FastRocks is more efficient than Rocks, it may be difficult to modify it for similar games.

0 1 2 3 4 5 6 7 8 9 10

0		W	L	W	L	W	L	W	L	W	L
1	W						_	_	_	_	W
2	L		L				_			W	
3	W	W	W	W	W	W	W	W	W	W	W
4	L	W	L	W	L	W	L	W	L	W	L
5	W	W	W	W	W	W	W	W	W	W	W
6	L	W	L	W	L	W	L	W	L	W	L
7	W	W	W	W	W	W	W	W	W	W	W
8	L	W	L	W	L	W	L	W	L	W	L
9	W	W	W	W	W	W	W	W	W	W	W
0	L	W	L	W	L	W	L	W	L	W	L