

Dynamic Programming for Boolean decisions

Lecture 07.03

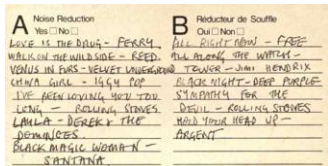
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Mix tapes

Subset sum

Mixtapes

- ❑ The *mixtape problem* is inspired by making [musical mixes](#) on cassettes (and later CDs)
- ❑ Given a set of songs with their durations, the question is whether these songs can be divided into 2 subsets where the total duration of each subset is the same



Mixtape problem

Input. The durations of n songs d_1, d_2, \dots, d_n in minutes (integers).

Output. Yes, if the songs can be divided into two groups, such that each group has the same total duration.

No, otherwise.

Sample problem instance

	1	2	3	4	5	6	7
min	3	2	3	2	2	5	3

A total of 7 songs given with their durations.

The output for this instance of the problem is 'yes'. This is because the songs can be divided into two groups that both have a total duration of 10 min.

3	2	3	2	2	5	3
1	2	3	4	5	6	7

The blue and red songs have the same total duration of 10 min.

DP solution: brainstorming

- ❑ What would help us to know if a set of numbers can be divided into 2 subsets with equal sums?
- ❑ How can we find out if there is a subset with a given sum?
- ❑ What are optimal subproblems?

Think of an optimal solution to a subset sum

- If there is a subset with total duration D , and it contains song i , then there also should be a subset with duration $D-d_i$



- As always, we can start by checking if all possible durations from 1 to D can be obtained from a current set, and we will reuse this knowledge to obtain an answer for duration D

Example

3	2	1	4	1	5
1	2	3	4	5	6

- ❑ First, we compute total duration:

$$3+2+1+4+1+5 = 16$$

- ❑ The task becomes to find out if there is a subset that sums up to $16/2 = 8$
- ❑ We will try methodically to fit each song into the solution, checking if the following total durations are possible: 0,1,2,3,4,5,6,7 and finally 8.
- ❑ The check will produce a boolean value: Y(True) or N(False)

Game of Rocks

Optimal game strategy

Game: 1-2 rocks

- 2 players
- 2 piles of rocks:

with n and m rocks respectively

- Each turn, one player may take either 1 rock (from either pile) or 2 rocks (one from each pile)
- Once the rocks are taken, they are removed from play
- The player that takes the last rock wins



Winning strategy with DP

- ❑ To find the winning strategy for the $m + n$ game, we first construct an $m \times n$ table R .
- ❑ If Player 1 can always win the $n + m$ game, then we would say $R(n, m) = W$, but if Player 1 has no winning strategy against a player that always makes the right moves, we would write $R(n, m) = L$.
- ❑ Computing $R(n, m)$ for arbitrary n and m seems difficult, but we can build on smaller values.

Fill DP table with game outcomes

- We can proceed filling in R in this way by noticing that for the entry (i, j) to be L , all the entries above, diagonally to the left, and directly to the left, must be W .

- These entries:

$((i - 1, j), (i - 1, j - 1), (i, j - 1))$
correspond to the three possible moves that Player 1 can make.

	0	1	2	3	4	5	6	7	8	9	10
0		W	L	W	L	W	L	W	L	W	L
1	W	W	W	W	W	W	W	W	W	W	W
2	L	W	L	W	L	W	L	W	L	W	L
3	W	W	W	W	W	W	W	W	W	W	W
4	L	W	L	W	L	W	L	W	L	W	L
5	W	W	W	W	W	W	W	W	W	W	W
6	L	W	L	W	L	W	L	W	L	W	L
7	W	W	W	W	W	W	W	W	W	W	W
8	L	W	L	W	L	W	L	W	L	W	L
9	W	W	W	W	W	W	W	W	W	W	W
10	L	W	L	W	L	W	L	W	L	W	L

Rocks: winning strategy

- ❑ The *Rocks* algorithm determines if Player 1 wins or loses.
- ❑ If Player 1 wins in an $n+m$ game, *Rocks* returns W. If Player 1 loses, *Rocks* returns L.
- ❑ We introduce an artificial initial condition, $R(0, 0) = L$ to simplify the pseudocode.

Algorithm *Rocks*(n, m)

$R[0, 0] \leftarrow L$

for i from 1 to n : # initialize rows

 if $R[i - 1, 0] = W$:

$R[i, 0] \leftarrow L$

 else:

$R[i, 0] \leftarrow W$

for j from 1 to m : # initialize columns

 if $R[0, j - 1] = W$:

$R[0, j] \leftarrow L$

 else:

$R[0, j] \leftarrow W$

for i from 1 to n :

 for j from 1 to m : # fill DP table

 if $R[i - 1, j - 1] = W$ and $R[i, j - 1] = W$ and $R[i - 1, j] = W$:

$R[i, j] \leftarrow L$

 else:

$R[i, j] \leftarrow W$

return $R[n, m]$

Using DP table for best strategy or game AI

- We can use the DP table to always play the winning strategy.
- If $R(n,m) = W$, and Player 1 starts first, he can always win: by taking the number of rocks which lead to the losing position of our opponent.
- If $R(n,m) = L$, then Player 1 can only hope that Player 2 does not use the same table, and makes a mistake.

	0	1	2	3	4	5
0		W	L	W	L	W
1	W	W	W	W	W	W
2	L	W	L	W	L	W
3	W	W	W	W	W	W
4	L	W	L	W	L	W
5	W	W	W	W	W	W

Winning strategy: example

	0	1	2	3	4	5
0		W	L	W	L	W
1	W	W	W	W	W	W
2	L	W	L	W	L	W
3	W	W	W	W	W	W
4	L	W	L	W	L	W
5	W	W	W	W	W	W

Player 1 takes (1,1).

Winning strategy example

	0	1	2	3	4	5
0		W	L	W	L	
1	W	W	W	W	W	
2	L	W	L	W	L	
3	W	W	W	W	W	
4	L	W	L	W	L	
5						

Player 1 takes (1,1).

No matter what Player 2 does, it leads to the winning state of Player 1.

Say, Player 2 takes (1,0)

Winning strategy example

	0	1	2	3	4	5
0		W	L	W	L	
1	W	W	W	W	W	
2	L	W	L	W	L	
3	W	W	W	W	W	
4						
5						

Player 2 takes (1,0)

Winning strategy example

	0	1	2	3	4	5
0		W	L	W	L	
1	W	W	W	W	W	
2	L	W	L	W	L	
3	W	W	W	W	W	
4						
5						

Player 1 should take (1,0).

Winning strategy example

	0	1	2	3	4	5
0		W	L	W	L	
1	W	W	W	W	W	
2	L	W	L	W	L	
3						
4						
5						

Player 1 takes (1,0).

Winning strategy example

	0	1	2	3	4	5
0		W	L	W	L	
1	W	W	W	W	W	
2	L	W	L	W	L	
3						
4						
5						

Player 2 takes (1,1).

Winning strategy example

	0	1	2	3	4	5
0		W	L	W		
1	W	W	W	W		
2						
3						
4						
5						

Player 2 takes (1,1).

Winning strategy example

	0	1	2	3	4	5
0		W	L	W		
1	W	W	W	W		
2						
3						
4						
5						

Player 1 should take (1,1).

Winning strategy example

	0	1	2	3	4	5
0		W	L			
1						
2						
3						
4						
5						

Player 1 takes (1,1).

At this point the victory for Player 1 is guaranteed.

Identifying patterns

- A faster algorithm relies on the simple pattern in R, and checks if n and m are both even, in which case the player 1 loses.
- However, though *FastRocks* is more efficient than *Rocks*, it may be difficult to modify it for similar games.

	0	1	2	3	4	5	6	7	8	9	10
0		W	L	W	L	W	L	W	L	W	L
1	W	W	W	W	W	W	W	W	W	W	W
2	L	W	L	W	L	W	L	W	L	W	L
3	W	W	W	W	W	W	W	W	W	W	W
4	L	W	L	W	L	W	L	W	L	W	L
5	W	W	W	W	W	W	W	W	W	W	W
6	L	W	L	W	L	W	L	W	L	W	L
7	W	W	W	W	W	W	W	W	W	W	W
8	L	W	L	W	L	W	L	W	L	W	L
9	W	W	W	W	W	W	W	W	W	W	W
10	L	W	L	W	L	W	L	W	L	W	L